



Service Restorability in Degree-Based Wavelength Division Multiplexing Networks

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Problem

- Internet is a complex, distributed system
- Better design strategies are required for reliable, high confidence, distributed systems
- Service restorability is one network-wide parameter and depends on the physical and virtual topology
- Previous network topology generator models were structural, but real-world networks follow degree-based models
- How do degree-based models affect restorability?

Graph-Based Topology Models

- Random
- Regular
- Well-known
- Structural/Hierarchical
- Degree-based/Power law

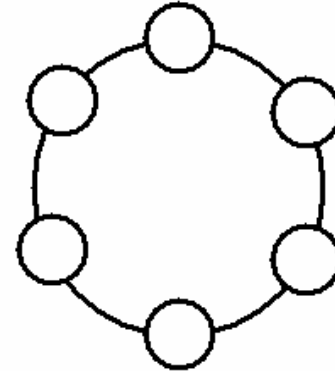
Flat Random Model

Method	Edge Probability	Reference
Pure Random	p	-
Waxman	$\alpha \exp(-d/\beta L)$ $d = \text{distance}$ $L = \text{max. distance}$ $\alpha > 0, \beta \leq 1$	Waxman 1988
Exponential	$\alpha \exp(-d/(L-d))$ $d = \text{distance}$ $L = \text{max. distance}$ $\alpha > 0$	Zegura <i>et al.</i> 1997
Locality	$\alpha, \text{ if } d < r$ $\beta, \text{ if } d \geq r$	Zegura <i>et al.</i> 1997

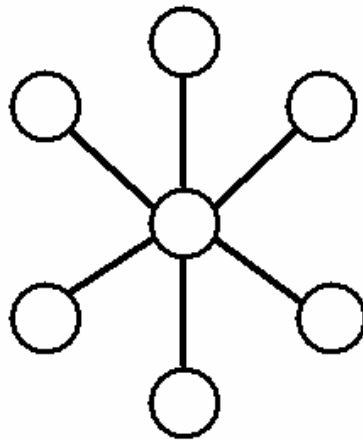
Regular Topologies



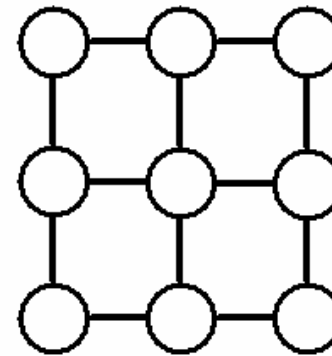
Linear chain



Ring

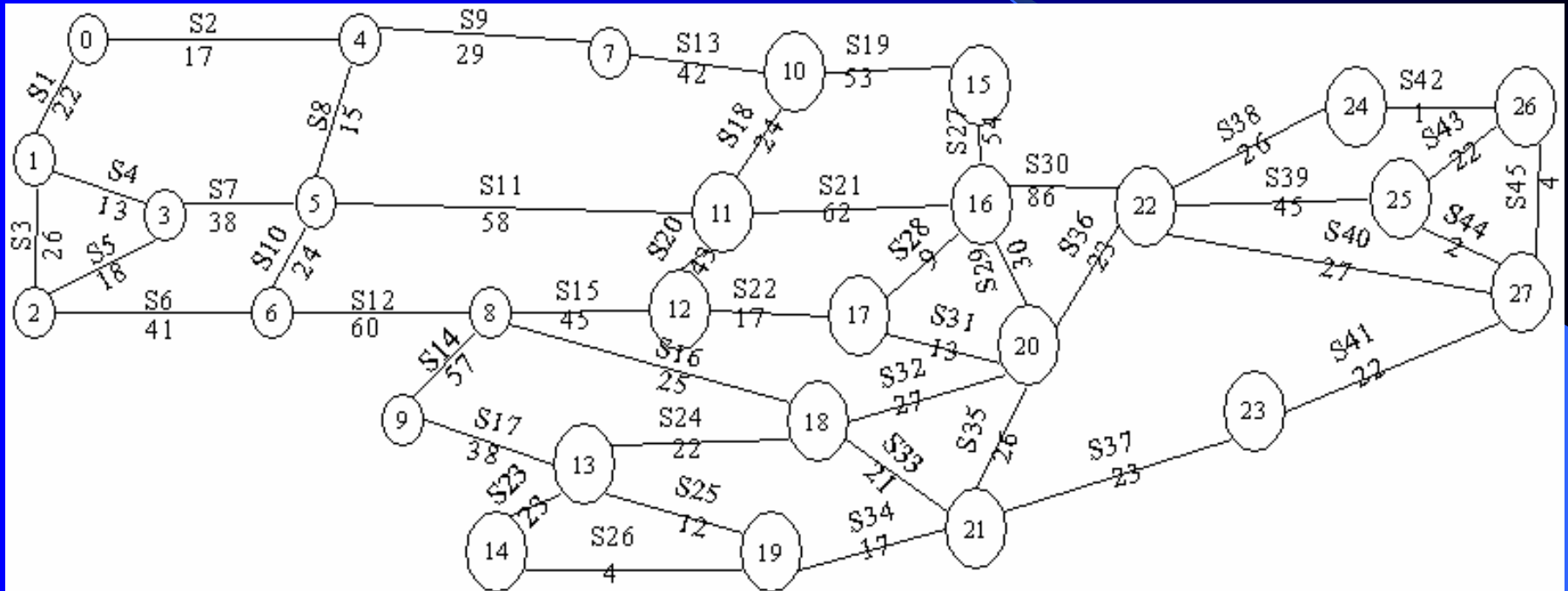


Star



Mesh

Well-known Topology NSFnet



Structural/Hierarchical Models

- N-Level (Zegura, Calvert, and Donahoo 1997)
 - Iteratively expand individual nodes into graphs
 - Edges of original graph attached to nodes in replacement graphs
- Transit-Stub (Zegura, Calvert, and Donahoo 1997)
 - Generate connected graph; each node represents a transit domain.
 - Each node replaced by another connected graph, representing backbone topology of one transit domain.
 - For each node in each transit domain, generate connected random graphs representing stub domains.
 - Add “extra” connectivity controlled by parameters.
- Structural models popular for building large networks with low average node degree
- Good attempt to mimic the structure of the Internet

Degree-Based/Power Law Model

- Faloutsos *et al.* 1999

$y \propto x^a$
 y and x are metrics
 a is a constant
 \propto : "proportional to"

- Power-Law 2 (outdegree exponent): The frequency, f_d , of an outdegree, d , is proportional to the outdegree to the power of a constant, O :

$f_d \propto d^O$
 f_d : frequency of outdegree, d
 O : constant
 \propto : "proportional to"

- $O = -2.2$ for Internet interdomain *and* router graphs
- Tangmunarunkit *et al.* 2002 compared structural and degree-based models and found that degree-based models model the Internet better even though they make no assumptions about hierarchical structure!

Degree-Based Barabasi model

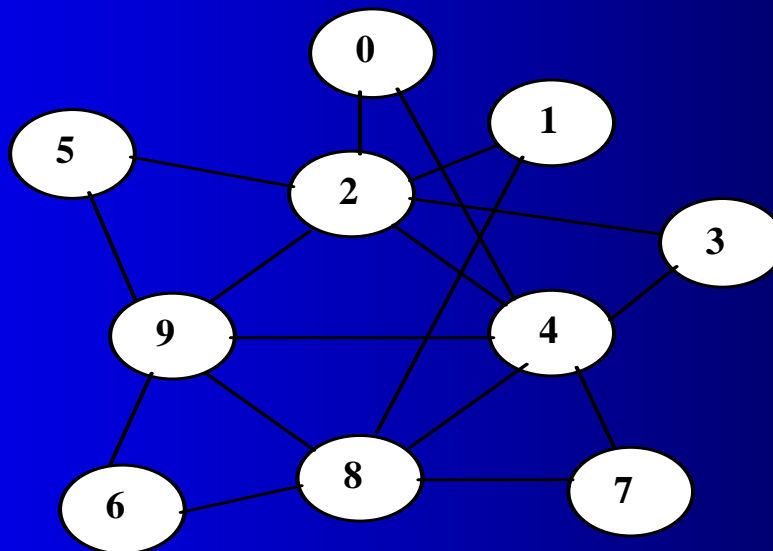
- Barabasi and Albert 1999
- Generates topologies based on two parameters
 - Incremental growth
 - Add new vertex with m edges
 - Preferential connectivity

$$\Pi(k_i) = k_i / \sum_j k_j$$

k_i = connectivity of vertex i

- Exponent of resulting networking approx. 2.9
- *BRITE* software implements this model

10-node Barabasi Model Graph



Degree-Based Power-Law Out-Degree (PLOD) Method

- Palmer and Steffan 2000
- Generates topologies based on out-degrees of nodes
- Assigns out-degree numbers to each node based on the expression

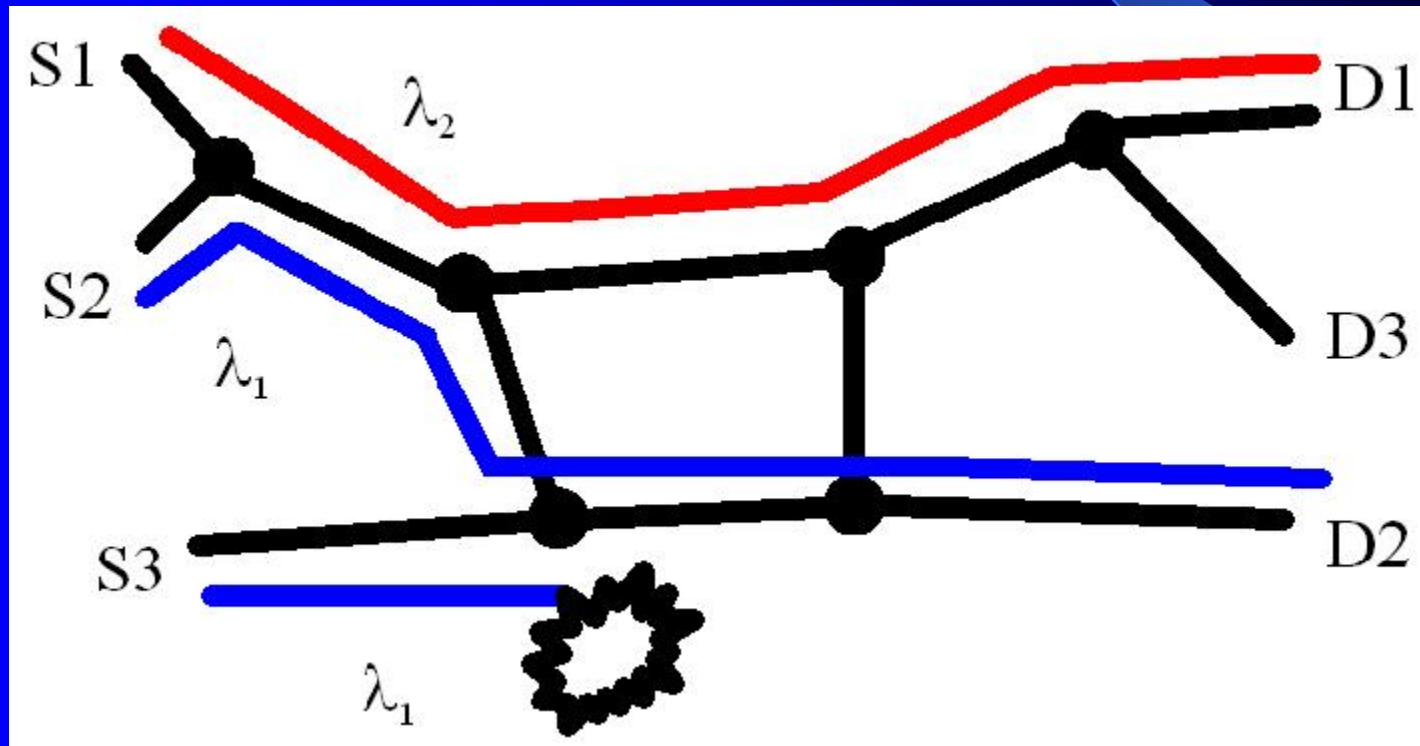
$$\beta x_i^{-\alpha}$$

- Then edges are generated according to the edge placement strategy
- α and β should be carefully adjusted to get the desired number of edges
- *Nem: Network manipulator* software

PLOD cont.

- Edge placement strategy:
 - Two nodes are picked randomly.
 - An edge is generated between these two nodes if there is no edge between them and if their out-degree numbers are not zeros; decrease the out-degree number by 1 if an edge is placed.
 - If the two nodes picked already have an edge between them or if their out-degree numbers are zeros then the process is repeated by picking two more nodes randomly.
 - This is done until all the nodes are assigned edges.

Wavelength Division Multiplexing (WDM)



Service Restorability

- Service Restorability is a network-wide parameter
- Restorability depends on the physical and virtual topology of the network
- The definition of restorability is the average fraction of failed working capacity that is restored for a specific failure scenario

Service Restorability in WDM Networks

- PRSA gives percentage of wavelengths assigned for given traffic.
- In this work, restorability is defined as percentage of the requested wavelengths assigned by PRSA averaged over a specific failure scenario.
- A failure scenario can be single-link failure or dual-link failure.
- A link is a single optical fiber.

Parallel Recombinative Simulated Annealing (PRSA) Algorithm

Initialize the Temperature (SA)

Initialize population with n chromosomes (GA)

Repeat for *max* generations

Do $n/2$ times

- Select 2 parent chromosomes at random (GA)
- Generate 2 children using crossover and mutation (GA)
- Hold competitions using the *Metropolis criterion* between children and parents (SA)
- Overwrite parents with trial winner

Lower the Temperature (SA)

Send/Receive migrants to/from other processors

Objective

- To find service restorability in degree-based wavelength division multiplexing networks

Methodology

- Single-link and dual-link failure simulations are performed to analyze restorability
- Networks with size 10, 15, 20, 25, and 50 nodes are considered for simulations
- Barabasi model graphs were generated with *BRITE* software
- PLOD model graphs were generated with *nem* software

Methodology

Steps involved:

For each topology generated

Convert the topology file in BRITE/nem format to the format used by PRSA.

Generate random traffic for this topology.

Keep scaling the traffic until 100% of wavelengths are assigned by PRSA.

Once the traffic is scaled, delete one link at a time for single link failures (delete two links at a time for dual link failures) and run the PRSA algorithm.

Find the average fraction of the failed working capacity that is restored over all single link failure scenarios (dual link failure scenarios). This is the restorability parameter for single-link failures (dual-link failures).

End

Single-link Failure Restorability

<i>Number of Nodes</i>	<i>Number of Links</i>	<i>Average</i>		<i>Standard Deviation</i>	
		<i>Barabasi</i>	<i>PLOD</i>	<i>Barabasi</i>	<i>PLOD</i>
10	17	98.78	98.01	1.32	4.83
15	27	98.83	98.74	2.11	1.96
20	37	99.82	99.32	0.49	1.46
25	47	98.63	97.45	1.66	2.25
50	97	99.50	99.93	0.80	0.48

Dual-link Failure Restorability

<i>Number of Nodes</i>	<i>Number of Links</i>	<i>Average</i>		<i>Standard Deviation</i>	
		<i>Barabasi</i>	<i>PLOD</i>	<i>Barabasi</i>	<i>PLOD</i>
10	17	97.32	97.05	4.58	6.53
15	27	96.38	98.16	3.34	2.88
20	37	98.91	98.32	1.88	2.11
25	47	97.72	95.83	2.39	2.88

Random vs. Degree-Based Models

- A 25-node Waxman's model is compared with degree-based models

<i>Model</i>	<i>Average</i>	<i>Standard Deviation</i>
Waxman's	98.42	2.15
Barabasi	98.63	1.66
PLOD	97.45	2.25

- Averages are very similar for all the models
- All the models have high restorability under this metric

Execution Times

Number of nodes	Time taken for each execution by PRSA	Number of executions	
		Single-link	Dual-link
10	40secs	17	136
15	2mins	27	351
20	6mins	37	666
25	12mins	47	1081
50	2hrs	97	4656

Conclusions

- Results show that degree-based models have high restorability even with dual-link failures

Future Work

- Need to extend to a set of networks for a particular number of nodes
- Need to extend to larger networks, which requires sampling
- Comparisons with structural models over larger range
- Study of node failures because degree-based networks have a few nodes with high node degree

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